

**Example-** An engineering problem involving logical conditions with 4 inputs results in the Truth Table shown where “X” means we don’t care about the output under these conditions (e.g., maybe this combination of inputs is physically impossible).

<i>a</i>	<i>b</i>	<i>c</i>	<i>d</i>	<i>F</i>
0	0	0	0	<b>0</b>
0	0	0	1	<b>0</b>
0	0	1	0	<b>0</b>
0	0	1	1	<b>1</b>
0	1	0	0	<b>1</b>
0	1	0	1	<b>X</b>
0	1	1	0	<b>X</b>
0	1	1	1	<b>1</b>
1	0	0	0	<b>0</b>
1	0	0	1	<b>0</b>
1	0	1	0	<b>X</b>
1	0	1	1	<b>1</b>
1	1	0	0	<b>0</b>
1	1	0	1	<b>1</b>
1	1	1	0	<b>X</b>
1	1	1	1	<b>X</b>

We will use a four variable K-Map to determine Boolean functions in the simplified sum-of-products and product-of-sums forms for this problem.

- First create K-Map directly from Truth Table.

<i>ab\cd</i>	<b>00</b>	<b>01</b>	<b>11</b>	<b>10</b>
<b>00</b>	<b>0</b>	<b>0</b>	<b>1</b>	<b>0</b>
<b>01</b>	<b>1</b>	<b>X</b>	<b>1</b>	<b>X</b>
<b>11</b>	<b>0</b>	<b>1</b>	<b>X</b>	<b>X</b>
<b>10</b>	<b>0</b>	<b>0</b>	<b>1</b>	<b>X</b>

- Next, we will use this K-Map to directly express  $F$  in a simplified sum-of-products form. The following sequence allows the coverage of all squares/minterms in groups of four by taking advantage of the don't-care squares/minterms.

$ab \backslash cd$	00	01	11	10
00	0	0	1	0
01	1	X	1	X
11	0	1	X	X
10	0	0	1	X

$ab \backslash cd$	00	01	11	10
00	0	0	1	0
01	1	X	1	X
11	0	1	X	X
10	0	0	1	X

$ab \backslash cd$	00	01	11	10
00	0	0	1	0
01	1	X	1	X
11	0	1	X	X
10	0	0	1	X

The K-Map with the selected minterms filled-in.

$ab \backslash cd$	00	01	11	10
00	$m_0$	$m_4$	$m_3 = a'b'cd$	$m_2$
01	$m_4 = a'b'cd'$	$m_5 = a'b'cd$	$m_7 = a'bcd$	$m_6 = a'bcd'$
11	$m_{12}$	$m_{13} = abc'd$	$m_{15} = abcd$	$m_{14}$
10	$m_8$	$m_9$	$m_{11} = ab'cd$	$m_{10}$

The simplified sum-of-products form is:  $\underline{F_{sop} = a'b + bd + cd}$  ;  
a result with **three** terms and **six** literals.

- What if all the don't-care minterms/squares had been made "0"s?

<i>ab\cd</i>	00	01	11	10
00	0	0	1	0
01	1	0	1	0
11	0	1	0	0
10	0	0	1	0

The following sequence allows the best coverage of all squares/minterms under this scenario.

<i>ab\cd</i>	00	01	11	10
00	0	0	1	0
01	1	0	1	0
11	0	1	0	0
10	0	0	1	0

<i>ab\cd</i>	00	01	11	10
00	0	0	1	0
01	1	0	1	0
11	0	1	0	0
10	0	0	1	0

<i>ab\cd</i>	00	01	11	10
00	0	0	1	0
01	1	0	1	0
11	0	1	0	0
10	0	0	1	0

The simplified sum-of-products form is  $a'cd + b'cd + a'bc'd' + abc'd$  ; a result with **four** terms and **14** literals. Not nearly as compact.

- Next, we will use this K-Map to express  $F$  in a simplified product-of-sums form(s). To get  $F'$  in simplified sum-of-products form(s), the following sequences allow the coverage of all the “0” squares/minterms in groups of two or four by taking advantage of the don't-care squares/minterms.

$ab \backslash cd$	00	01	11	10
00	0	0	1	0
01	1	X	1	X
11	0	1	X	X
10	0	0	1	X

$ab \backslash cd$	00	01	11	10
00	0	0	1	0
01	1	X	1	X
11	0	1	X	X
10	0	0	1	X

and

$ab \backslash cd$	00	01	11	10
00	0	0	1	0
01	1	X	1	X
11	0	1	X	X
10	0	0	1	X

or

$ab \backslash cd$	00	01	11	10
00	0	0	1	0
01	1	X	1	X
11	0	1	X	X
10	0	0	1	X

Here,  $F' = b'c' + cd' + ac'd'$  or  $F' = b'c' + cd' + abd'$ .

Now, we take the complement of  $F'$  to express  $F$  in a simplified product-of-sums form by using DeMorgan's Theorem (i.e., take dual and then complement each literal).

$$\begin{aligned} F_{\text{pos1}} &= (F')' = (b'c' + cd' + ac'd')' \\ &= \underline{(b+c)(c'+d)(a'+c+d)} \end{aligned}$$

a result with **three** terms and **seven** literals (two complemented). Or,

$$\begin{aligned} F_{\text{pos2}} &= (F')' = (b'c' + cd' + abd')' \\ &= \underline{(b+c)(c'+d)(a'+b'+d)} \end{aligned}$$

a result with **three** terms and **seven** literals (three complemented). By one complement/NOT operation, the first option is slightly simpler.

- Finally, check these results using Truth Tables to show they agree for the "0"s and "1"s. We don't-care about the "X"s!

$$\underline{F_{\text{sop}} = a'b + bd + cd}$$

$a$	$b$	$c$	$d$	$a'$	$a'b$	$bd$	$cd$	$F_{\text{sop}}$	$F$
0	0	0	0	1	0	0	0	0	0
0	0	0	1	1	0	0	0	0	0
0	0	1	0	1	0	0	0	0	0
0	0	1	1	1	0	0	1	1	1
0	1	0	0	1	1	0	0	1	1
0	1	0	1	1	1	1	0	1	X
0	1	1	0	1	1	0	0	1	X
0	1	1	1	1	1	1	1	1	1
1	0	0	0	0	0	0	0	0	0
1	0	0	1	0	0	0	0	0	0
1	0	1	0	0	0	0	0	0	X
1	0	1	1	0	0	0	1	1	1
1	1	0	0	0	0	0	0	0	0
1	1	0	1	0	0	1	0	1	1
1	1	1	0	0	0	0	0	0	X
1	1	1	1	0	0	1	1	1	X

$$F_{\text{pos1}} = (b + c)(c' + d)(a' + c + d)$$

<i>a</i>	<i>b</i>	<i>c</i>	<i>d</i>	<i>a'</i>	<i>c'</i>	<i>b+c</i>	<i>c'+d</i>	<i>a'+c+d</i>	<i>F<sub>pos1</sub></i>	<i>F</i>
0	0	0	0	1	1	0	1	1	0	0
0	0	0	1	1	1	0	1	1	0	0
0	0	1	0	1	0	1	0	1	0	0
0	0	1	1	1	0	1	1	1	1	1
0	1	0	0	1	1	1	1	1	1	1
0	1	0	1	1	1	1	1	1	1	X
0	1	1	0	1	0	1	0	1	0	X
0	1	1	1	1	0	1	1	1	1	1
1	0	0	0	0	1	0	1	0	0	0
1	0	0	1	0	1	0	1	1	0	0
1	0	1	0	0	0	1	0	1	0	X
1	0	1	1	0	0	1	1	1	1	1
1	1	0	0	0	1	1	1	0	0	0
1	1	0	1	0	1	1	1	1	1	1
1	1	1	0	0	0	1	0	1	0	X
1	1	1	1	0	0	1	1	1	1	X

$$F_{\text{pos2}} = (b + c)(c' + d)(a' + b' + d)$$

<i>a</i>	<i>b</i>	<i>c</i>	<i>d</i>	<i>a'</i>	<i>b'</i>	<i>c'</i>	<i>b+c</i>	<i>c'+d</i>	<i>a'+b'+d</i>	<i>F<sub>pos2</sub></i>	<i>F</i>
0	0	0	0	1	1	1	0	1	1	0	0
0	0	0	1	1	1	1	0	1	1	0	0
0	0	1	0	1	1	0	1	0	1	0	0
0	0	1	1	1	1	0	1	1	1	1	1
0	1	0	0	1	0	1	1	1	1	1	1
0	1	0	1	1	0	1	1	1	1	1	X
0	1	1	0	1	0	0	1	0	1	0	X
0	1	1	1	1	0	0	1	1	1	1	1
1	0	0	0	0	1	1	0	1	1	0	0
1	0	0	1	0	1	1	0	1	1	0	0
1	0	1	0	0	1	0	1	0	1	0	X
1	0	1	1	0	1	0	1	1	1	1	1
1	1	0	0	0	0	1	1	1	0	0	0
1	1	0	1	0	0	1	1	1	1	1	1
1	1	1	0	0	0	0	1	0	0	0	X
1	1	1	1	0	0	0	1	1	1	1	X